## MSc Geomatics thesis presentation

Validation and automatic repair of planar
partitions using a constrained triangulation
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## TUDelft

## Representing thematic information

- Pierre Charles Dupin's 1819 map
- Subdividing the feature space of that theme
- Distinct visual representation for each class



# Digital thematic maps 

- CGIS, SYMAP (Late ‘60s)
- Polygons to represent boundaries
- Well suited for a computer




## Why planar partitions?

- Simple constraints: no gaps, no overlaps, (no disjoint regions)
- Easy to answer questions:
- (Aggregation) What is the total area of the features of type A and B ?
- (Topology) Are features A and B adjacent?


## CORINE E41N27



## The problem

- Validate a planar partition
- If it's invalid, automatically repair it.


## The problem

- Validate a planar partition
- If it's invalid, automatically repair it.


Why is it hard?


## Errors in computer operations

## Orientation test



Polygons with holes
Valid polygons I
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## Validation process

- Rings:
- Closed, not self-intersecting, not zero area, correct winding, ...
- Polygons:
- Nested rings, connected interior, not zero area, ...
- Planar Partitions:
- No gaps, no overlaps, no disjoint regions...

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## Polygon validation constraints

- ArcGIS: too short line segments, unclosed rings, self-intersections, incorrect ring ordering.
- JTS/GEOS: self-touching rings, zero area rings, zero area polygons, improperly nested rings, duplicate vertices, spikes and gores, touching parts, crossing rings.
- Oracle Spatial: polygon with fewer than 4 vertices, unclosed rings, selfintersections, touching rings, overlapping rings, points too close together, wrong orientation.


Problems are unavoidable

ArcGIS and the zero area polygon

## Validation of planar partitions

- Nearly impossible without topology: finding gaps
- Creation of a planar graph based representation
- Plümer and Groger: no dangling edges, no zero-length edges, planarity, no holes, no self-intersections, no overlaps, connectivity.




## Repair of planar partitions using snapping

## Thresholds

## Repair using snapping

- Extensively available: ArcGIS, FME, GRASS, Radius Topology
- Possibilities: point to point, point to line, line to line



Repair using snapping

## Low resolution data

High resolution data


Clean-up required

## After snapping: <br> constraints for planar partition repair

- Radius Topology: share node, node-split-edge, edge-split-edge
- GRASS: break at intersections, remove duplicate line segments, remove dangling edges, remove bridges, remove vertices within a threshold of a line segment, remove too small areas, remove too small angles
- But still, no guarantees


## Topological planar partition repair

- Use topological constraints instead
- Snapping can still be performed
- Available in ArcGIS: must not overlap, must not have gaps, must not overlap with, must not have dangles, must not have pseudonodes, must not self intersect





## Topological repair in ArcGIS



## Topological repair in ArcGIS

## Manual editing of topology

## The solution

- Repair individual polygons.
- Create a triangulation containing every edge of every polygon.
- Tag every triangle with the polygons it belongs to.
- Re-tag areas with multiple or no tags, according to predefined
 criteria.
- Reconstruct the polygons in the triangulation.


Triangulations


Triangulations
Polygons

The constrained Delaunay triangulation

- Delaunay triangulation: empty circle property, uniqueness
- Constrained edges



## Repair an individual polygon

- Use the same techniques devised to repair planar partitions.
- Create a triangulation from the polygon.
- Iteratively define exterior and interior when passing a constrained edge.

- Reconstruct polygon.


## Create triangulation

- Add every edge of every (now valid) polygon to the triangulation as a constrained edge.
- Track whenever constrained edges are split.


## Tag triangulation

- Mark each triangle with the polygons that it belongs to.
- No tags = gap
- Multiple tags = overlap



## Repair operations

- At the end, ensure that each triangle has exactly one tag.
- Some possible options:
- Assign triangle or region to the neighbour present on most sides.
- Assign triangle or region to the neighbour with the longest boundary.
- Assign region to the class with the highest priority.


Repair operations


Triangle with
longest boundary


Random region


Region with longest boundary

## Polygon reconstruction

- Recursively create a chain of edges representing all boundaries (and some connecting segments).
- Cut where more than two edges join.
- Join small chains in the correct order to form rings.



## The prototype

- C++ with CGAL and OGR
- Open source and freely available

00 ThesisFinal - Debugger Console

Triangles: 82661
Adding a new set of polygons to the triangulation...
File opened.
Name: /Volumes/Buffalo/corine/100KME40N32.shp
Type: ESRI Shapefile
Layers: 1
Layer[1]: 2081 polygons \{
string CODE_00
double
AREA
\}
(213) OB: Self intersecting.

Splitting ring (2169 nodes) at (4090897.81947201,3205131.85136284).
Created 2 rings.
Outer rings: 1 inner rings: 1
(295) OB: Self intersecting.

Splitting ring (171 nodes) at (4014338.28224871,3205959.57513961).. Created 2 rings
Outer rings: 1 inner rings: 1
(560) OB: Self intersecting.

Splitting ring ( 827 nodes) at (4006889.99296177,3214962.87908400).. Created 2 rings.
Outer rings: 1 inner rings: 1
(661) OB: Self intersecting.

Splitting ring (137 nodes) at (4052575.12529232,3220418.25137708).. Created 2 rings.
Outer rings: 1 inner rings: 1
(954) OB: Self intersecting.

Splitting ring (206 nodes) at (4041756.04611424,3231507.73413940)..
Created 2 rings.
Outer rings: 1 inner rings: 1
(983) OB: Self intersecting.

Splitting ring (135 nodes) at (4015515.72557781,3236851.22009638)..
Created 2 rings.
Outer rings: 1 inner rings: 1
(1219) OB: Self intersecting.

Splitting ring (186 nodes) at ( $4028234.13326965,3246289.96980117$ )..
Created 2 rings
Outer rings: 1 inner rings: 1
(1264) OB: Self intersecting.

Splitting ring (952 nodes) at (4041062.88515360,3236397.85638396).. Created 2 rings.

Debugging terminated.

## Comparisons with other software

- Test polygons with specific problems (ArcGIS)
- Significant differences in interpretation

- Standards specify how to define a certain polygon, not how to interpret an existing one



## Comparisons with other software

- Large "normal" data sets (ArcGIS, FME, GRASS)

| Test | CT <br> memory | time | ArcGIS <br> memory | time | FME <br> memory | time | memory | time |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| E41N27 | 124 MB | 46 s | 145 MB | 1 m 3 s | 158 MB | 31 s | 59 MB | 3 m 09 s |
| 4tiles | 100 MB | 3 m 25 s | 113 MB | 37 s | 105 MB | 31 s | 49 MB | 53 s |
| 16tiles | 1.51 GB | 1 h 20 m | crashes | - | 636 MB | 15 m 48 s | crashes | - |
| Mexico | 983 MB | 18 m 53 s | 216 MB | $>1 \mathrm{~d}$ | 264 MB | 2 m 45 s | 408 MB | 11 m 38 s |

- Good performance, considering that it does much more


## Conclusions

- Planar partition validation and automatic repair with a constrained triangulation is theoretically simple, yet powerful.
- It keeps a valid topology throughout, without a complex set of rules to check every step of the way.
- Changes that are made to the triangulation have only a local effect.
- New repair operations, based on different criteria, can be easily implemented without breaking the validity of the planar partition.
- Snapping is possible, but not required.


## Future work

- Optimisations for simpler polygons
- Improved algorithms for extracting polygons from a triangulation
- Eliminating memory limitations
- Improving the order of point insertions
- Extension to 3D
- Implementation in a database


## Questions?

http://www.gdmc.nl/~ken/thesis.pdf


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