Storing and analysing **MASSIVE** aerial LiDAR datasets in a DBMS

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- Computers have many problems dealing with billions of points:
 - Storing is OK
 - Visualisation is still a challenge
 - Processing and analysis are very problematic
- Processing operations:
 - derivation of slope/aspect,
 - conversion to grid format,
 - calculations of area/volumes,
 - viewshed analysis,
 - creation of simplified DTM,
 - extraction of bassins,
 - etc.

Advances in technologies to collect data are far superior to our ability to process data.



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LiDAR datasets are formed by scattered points in 3D space, which are the samples of a surface that can be projected on the horizontal plan.





Original LiDAR points



Raster representation



Raster representation



TIN (with Delaunay triangles)



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- Terrasolid: max is main memory
- ArcGIS: Terrain type (hierarchical structure)
- Oracle Spatial 11g: Point Cloud & TIN types
- External memory algorithms [AAD06, ADHZ06]
- Streaming of geometries of Isenburg *et al.* [ILSS06, ILS⁺06]

- **1** Storing independently triangles (\sim OGC)
- 2 Triangle-based data structure used by triangulation libraries [BDP⁺02]
- 3 Edge-based data structure (e.g. half-edge [M88])

- Goes beyond the usual "store points and edges/triangles"
- Ideas come from data structures for compression of graphs [BBCK05]



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Stars in a DBMS

ID	X	у	Z	star[]
1	3.21	5.23	2.11	[2,44,55,61,23]
2	5.19	29.01	4.55	[7,98,111,233,222]
3	22.43	15.99	8.19	[99,101,73,23]
5674	221.19	15.23	37.81	[309,802,793,1111]

Advantages:

- **1** Only **one** table with id x y z star
- 2 No spatial index needed: fetching of triangles based on "walking"
- 3 Star column need not be filled (\sim Simple Features)
- 4 Local updates are possible (insertion and removals)

Point Location = "Walking" in the triangulation



(Can be made efficient with some tricks [MSZ99])

Range Queries: also uses the triangulation





Figure from Martin Isenburg's presentation at GIScience 2006 [ILSS06]

Streaming of geometries to construct massive TINs



Experiments with AHN2 datasets



	# pts	# triangle	degree _{avg}	degree _{max}
20tiles	281 884 687	563 768 199	6.00	63
g37en1_15	8 605 090	17 201 289	6.00	39

Experiments with AHN2 datasets



Examples of queries: statistics about convex hull

20tiles=# select count(id) from points where is_convexhull(star) is true count

1173 (1 row) Time: 333050.861 ms



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1173 (1 row) Time: 333050.861 ms



20tiles=# select avg(degree(star)) from points; avg ------5.9999958142601850 (1 row) Time: 332265.041 ms





g37en1_15=# select degree(star) as degree, count(id) from points group by degree order by degree;

3 65620 4 844625 5 2277911 6 2484212 7 2005407 8 698540 9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	degree	count	
3 65620 4 844625 5 2277911 6 2484212 7 2005407 8 698540 9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	+		
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6 2484212 7 2005407 8 698540 9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	5	2277911	
7 2005407 8 698540 9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	6	2484212	
8 698540 9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	7	2005407	
9 170214 10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	8	698540	
10 37534 11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	9	170214	
11 9587 12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	10	37534	
12 3395 13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	11	9587	
13 1552 14 772 15 456 (truncated) 37 3 39 1 (35 rows)	12	3395	
14 772 15 456 (truncated) 37 3 39 1 (35 rows)	13	1552	
15 456 (truncated) 37 3 39 1 (35 rows)	14	772	
(truncated) 37 3 39 1 (35 rows)	15	456	
37 3 39 1 (35 rows)	(trur	ncated)	
39 1 (35 rows)	37	3	
(35 rows)	39	1	
	(35 rows)	1	
	Time 2010/	00 017	

Examples of queries: point location

g37en1_15=# select point_location(84111, 446666, 0, 0.2); WARNING: #of samples checked: 24 WARNING: start distance is 49.843357 WARNING: # of triangles visited is 222 point_location

(3672278,3695197,3695256) (1 row)

Time: 191.903 ms



Visualisation with a GIS





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Populating the database

	time pipepline	time B-tree	time tr	time lasblock	width _{max}
20tiles	178.3	24.65	36.8	crashed	53 003
g37en1_15	4.7	0.3	1.1	4.3	7 601